Today.





Secret Sharing.



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Correcting for loss or even corruption.

Share secret among *n* people.

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Two points make a line. Lots of lines go through one point.

A polynomial

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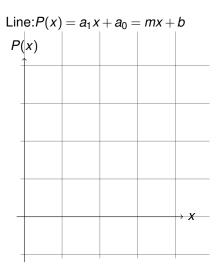
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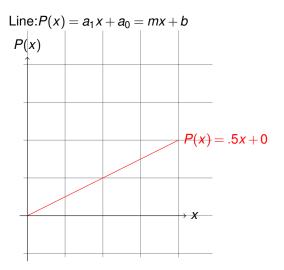
Polynomials P(x) with arithmetic modulo p: ¹ $a_i \in \{0, ..., p-1\}$ and

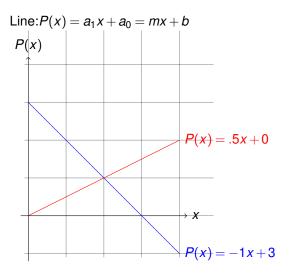
$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0 \pmod{p},$$
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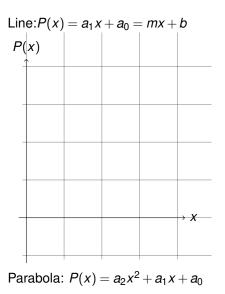
Line: $P(x) = a_1 x + a_0$

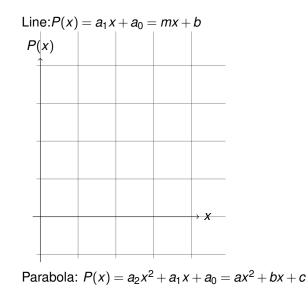
Line: $P(x) = a_1x + a_0 = mx + b$

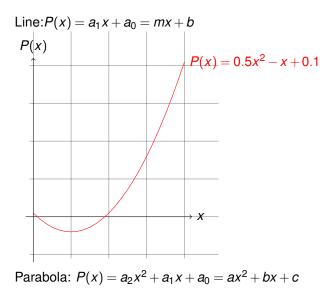


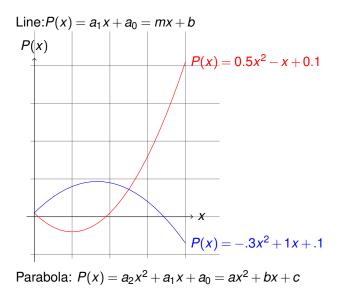


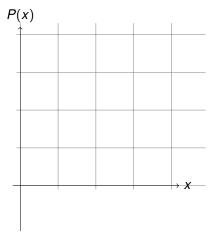


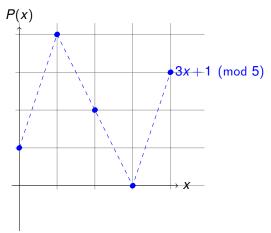


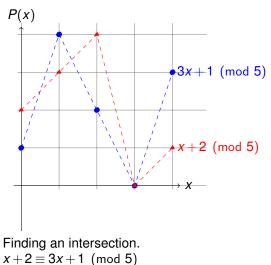




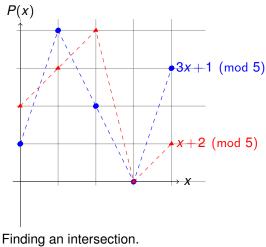




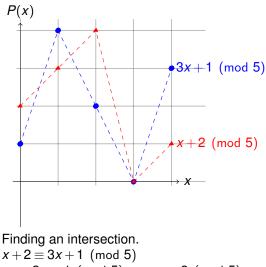




 $\implies 2x \equiv 1 \pmod{5}$



 $x + 2 \equiv 3x + 1 \pmod{5}$ $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5.



 $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Good when modulus is prime!! Two points make a line.

Fact: Exactly 1 degree $\leq d$ polynomial contains d + 1 points.²

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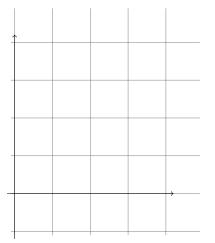
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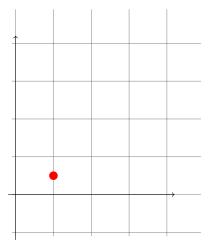
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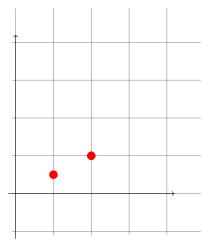
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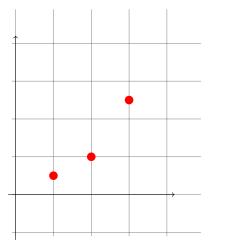
Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d + 1 pts.

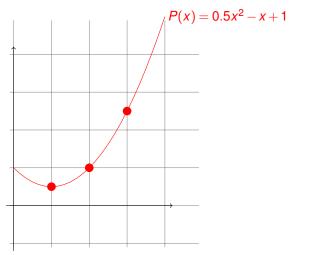
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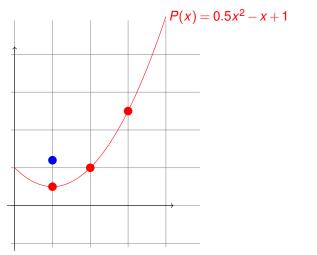


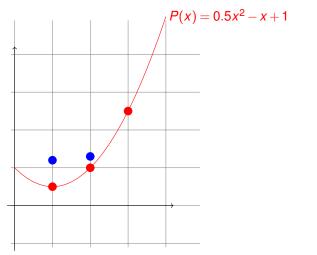


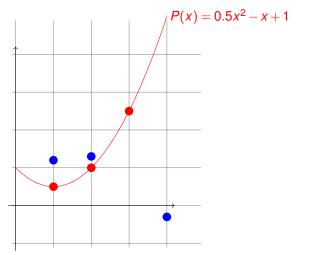


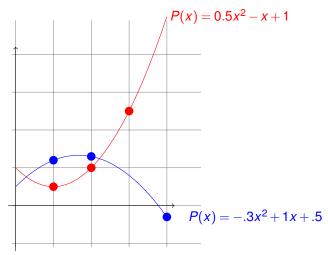




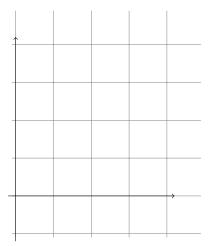


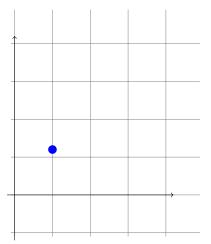


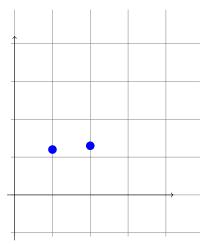


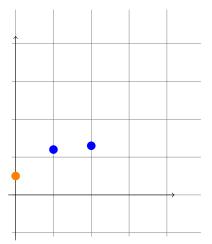


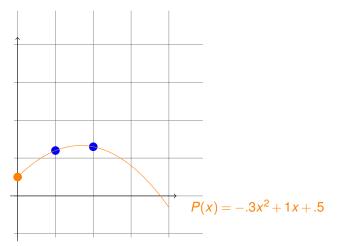
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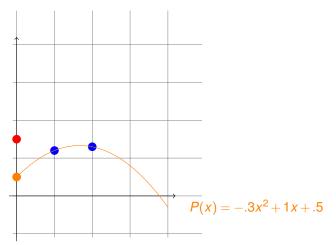


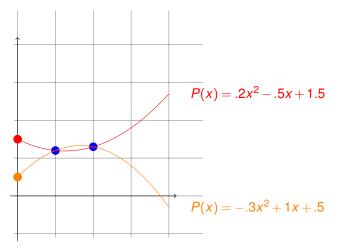


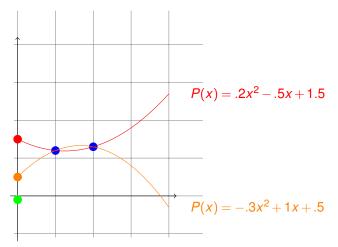


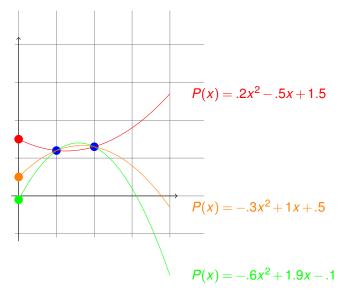


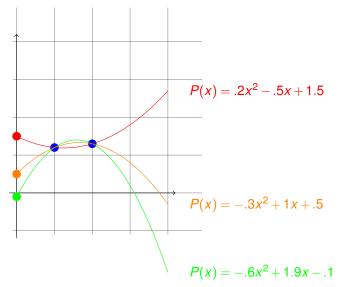












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From d + 1 points to degree d polynomial?

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 $x+2 \mod 5$.

For a quadratic polynomial, $a_2x^2 + a_1x + a_0$ hits (1,2); (2,4); (3,0).

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Multiplicative inverses due to gcd(x,p) = 1, forall $x \in \{1, ..., p-1\}$

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There exists a polynomial...

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Put the delta functions together.

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Assume two different polynomials Q(x) and P(x) hit the points.

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Assume two different polynomials Q(x) and P(x) hit the points.

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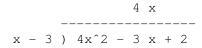
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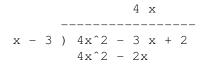
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Must prove Roots fact.





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$$4 x + 4 r 4$$

$$x - 3) 4x^{2} - 3 x + 2$$

$$4x^{2} - 2x$$

$$-----$$

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$$-----$$

$$4$$

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 $4x^2 - 3x + 2 \equiv (x - 3)(4x + 4) + 4 \pmod{5}$ In general, divide P(x) by (x - a) gives Q(x) and remainder r.

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In general, divide $P(x)$ by $(x - a)$ gives $Q(x)$ and remainder r .
That is, $P(x) = (x - a)Q(x) + r$

Lemma 1: P(x) has root *a* iff P(x)/(x-a) has remainder 0: P(x) = (x-a)Q(x).

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Roots fact: Any degree *d* polynomial has at most *d* roots.

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- Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.