1. Isoperimetric inequality for hypercube.

- 1. Isoperimetric inequality for hypercube.
- 2. Modular Arithmetic.

- 1. Isoperimetric inequality for hypercube.
- 2. Modular Arithmetic. Clock Math!!!

- 1. Isoperimetric inequality for hypercube.
- 2. Modular Arithmetic. Clock Math!!!
- 3. Inverses for Modular Arithmetic: Greatest Common Divisor.

- 1. Isoperimetric inequality for hypercube.
- 2. Modular Arithmetic. Clock Math!!!
- 3. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!

- 1. Isoperimetric inequality for hypercube.
- 2. Modular Arithmetic. Clock Math!!!
- 3. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!
- 4. Euclid's GCD Algorithm.

- 1. Isoperimetric inequality for hypercube.
- 2. Modular Arithmetic. Clock Math!!!
- 3. Inverses for Modular Arithmetic: Greatest Common Divisor. Division!!!
- 4. Euclid's GCD Algorithm. A little tricky here!

For 3-space:

For 3-space:

The sphere minimizes surface area to volume.

For 3-space:

The sphere minimizes surface area to volume. Surface Area: $4\pi r^2$, Volume: $\frac{4}{3}\pi r^3$.

For 3-space:

The sphere minimizes surface area to volume. Surface Area: $4\pi r^2$, Volume: $\frac{4}{3}\pi r^3$. Ratio: $1/3r = \Theta(V^{-1/3})$.

For 3-space:

The sphere minimizes surface area to volume.

Surface Area: $4\pi r^2$, Volume: $\frac{4}{3}\pi r^3$.

Ratio: $1/3r = \Theta(V^{-1/3})$.

Graphical Analog: Cut into two pieces and find ratio of edges/vertices on small side.

For 3-space:

The sphere minimizes surface area to volume.

Surface Area: $4\pi r^2$, Volume: $\frac{4}{3}\pi r^3$.

Ratio: $1/3r = \Theta(V^{-1/3})$.

Graphical Analog: Cut into two pieces and find ratio of edges/vertices on small side.

Tree: $\Theta(1/|V|)$.

For 3-space:

The sphere minimizes surface area to volume.

Surface Area: $4\pi r^2$, Volume: $\frac{4}{3}\pi r^3$.

Ratio: $1/3r = \Theta(V^{-1/3})$.

Graphical Analog: Cut into two pieces and find ratio of edges/vertices on small side.

Tree: $\Theta(1/|V|)$.

Hypercube: $\Theta(1)$.

For 3-space:

The sphere minimizes surface area to volume.

Surface Area: $4\pi r^2$, Volume: $\frac{4}{3}\pi r^3$.

Ratio: $1/3r = \Theta(V^{-1/3})$.

Graphical Analog: Cut into two pieces and find ratio of edges/vertices on small side.

Tree: $\Theta(1/|V|)$.

Hypercube: $\Theta(1)$.

Surface Area is roughly at least the volume!

Recursive Definition.

A 0-dimensional hypercube is a node labelled with the empty string of bits.

Recursive Definition.

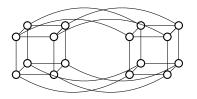
A 0-dimensional hypercube is a node labelled with the empty string of bits.

An *n*-dimensional hypercube consists of a 0-subcube (1-subcube) which is a n-1-dimensional hypercube with nodes labelled 0x (1x) with the additional edges (0x, 1x).

Recursive Definition.

A 0-dimensional hypercube is a node labelled with the empty string of bits.

An *n*-dimensional hypercube consists of a 0-subcube (1-subcube) which is a n-1-dimensional hypercube with nodes labelled 0x(1x) with the additional edges (0x, 1x).



Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S;

Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S; $|E \cap S \times (V - S)| \ge |S|$

Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S; $|E \cap S \times (V - S)| \ge |S|$

Terminology:

Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S; $|E \cap S \times (V - S)| \ge |S|$

Terminology: (S, V - S) is cut.

Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S; $|E \cap S \times (V - S)| \ge |S|$

Terminology: (S, V - S) is cut. $(E \cap S \times (V - S))$ - cut edges.

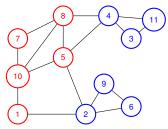
Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S; $|E \cap S \times (V - S)| \ge |S|$

Terminology: (S, V - S) is cut. $(E \cap S \times (V - S))$ - cut edges.

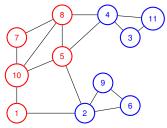
Thm: Any subset *S* of the hypercube where $|S| \le |V|/2$ has $\ge |S|$ edges connecting it to V - S; $|E \cap S \times (V - S)| \ge |S|$

Terminology: (S, V - S) is cut. $(E \cap S \times (V - S))$ - cut edges.

Restatement: for any cut in the hypercube, the number of cut edges is at least the size of the small side.

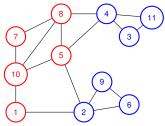


S is red, V - S is blue.



S is red, V - S is blue.

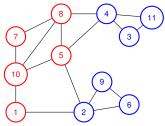
What is size of cut?



S is red, V - S is blue.

What is size of cut?

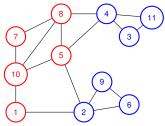
Number of edges between red and blue.



S is red, V - S is blue.

What is size of cut?

Number of edges between red and blue. 4.



S is red, V - S is blue.

What is size of cut?

Number of edges between red and blue. 4.

Hypercube: any cut that cuts off *x* nodes has $\ge x$ edges.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:**

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:** Base Case: n = 1

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:** Base Case: n = 1 V= {0,1}.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:**

Base Case: n = 1 V= {0,1}. S = {0} has one edge leaving.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:**

Base Case: n = 1 V= {0,1}. S = {0} has one edge leaving. $|S| = \phi$ has 0.

Proof of Large Cuts.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:**

Base Case: n = 1 V= {0,1}. S = {0} has one edge leaving. $|S| = \phi$ has 0.

Proof of Large Cuts.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side. **Proof:**

Base Case: n = 1 V= {0,1}. S = {0} has one edge leaving. $|S| = \phi$ has 0.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

Use recursive definition into two subcubes.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

Use recursive definition into two subcubes.

Two cubes connected by edges.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

Use recursive definition into two subcubes.

Two cubes connected by edges.

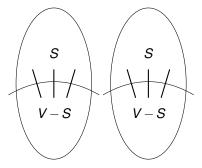
Case 1: Count edges inside subcube inductively.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

Use recursive definition into two subcubes.

Two cubes connected by edges.

Case 1: Count edges inside subcube inductively.



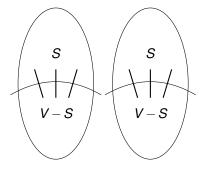
Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

Use recursive definition into two subcubes.

Two cubes connected by edges.

Case 1: Count edges inside subcube inductively.

Case 2: Count inside and across.

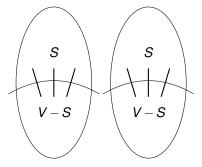


Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side.

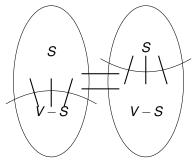
Use recursive definition into two subcubes.

Two cubes connected by edges.

Case 1: Count edges inside subcube inductively.



Case 2: Count inside and across.



Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them.

$$H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1),$ edges E_x that connect them.

$$H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2, |S_1| \le |V_1|/2$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them.

$$H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2, |S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2$, $|S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides. So by induction.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2$, $|S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides. So by induction. Edges cut in $H_0 \ge |S_0|$.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2, |S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides. So by induction. Edges cut in $H_0 \ge |S_0|$. Edges cut in $H_1 \ge |S_1|$.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2, |S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides. So by induction. Edges cut in $H_0 \ge |S_0|$. Edges cut in $H_1 \ge |S_1|$.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2, |S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides. So by induction. Edges cut in $H_0 \ge |S_0|$. Edges cut in $H_1 \ge |S_1|$.

 $\label{eq:constraint} \text{Total cut edges} \geq |\textbf{\textit{S}}_0| + |\textbf{\textit{S}}_1| = |\textbf{\textit{S}}|.$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step.

Recursive definition:

 $H_0 = (V_0, E_0), H_1 = (V_1, E_1)$, edges E_x that connect them. $H = (V_0 \cup V_1, E_0 \cup E_1 \cup E_x)$

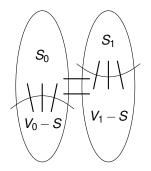
 $S = S_0 \cup S_1$ where S_0 in first, and S_1 in other.

Case 1: $|S_0| \le |V_0|/2, |S_1| \le |V_1|/2$ Both S_0 and S_1 are small sides. So by induction. Edges cut in $H_0 \ge |S_0|$. Edges cut in $H_1 \ge |S_1|$.

 $\text{Total cut edges} \geq |\mathcal{S}_0| + |\mathcal{S}_1| = |\mathcal{S}|.$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|. **Proof: Induction Step. Case 2.**

$$|S_0| \ge |V_0|/2.$$

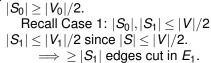


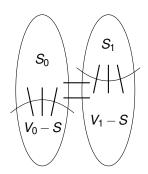
Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|. **Proof: Induction Step. Case 2.**

 S_0 $V_0 - S$ $V_1 - S$ $|S_0| \ge |V_0|/2.$ Recall Case 1: $|S_0|, |S_1| \le |V|/2$ $|S_1| \le |V_1|/2$ since $|S| \le |V|/2.$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

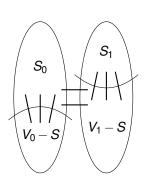
Proof: Induction Step. Case 2.





Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

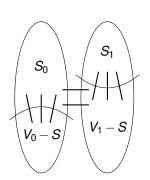
Proof: Induction Step. Case 2.



$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| &\leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| &\leq |V|/2. \\ &\implies &\geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| &\leq |V_0|/2 \end{split}$$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

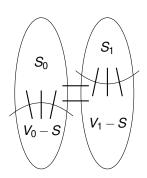
Proof: Induction Step. Case 2.



$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| \leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| \leq |V|/2. \\ &\implies \geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| \leq |V_0|/2 \\ &\implies \geq |V_0| - |S_0| \text{ edges cut in } E_0. \end{split}$$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.

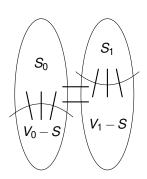


$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| \leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| \leq |V|/2. \\ &\implies \geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| \leq |V_0|/2 \\ &\implies \geq |V_0| - |S_0| \text{ edges cut in } E_0. \end{split}$$

Edges in E_x connect corresponding nodes.

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.

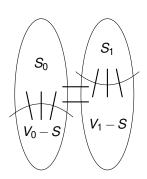


$$\begin{array}{l} \mathcal{S}_{0}|\geq |V_{0}|/2.\\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2\\ \mathcal{S}_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2.\\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}.\\ |S_{0}|\geq |V_{0}|/2\implies |V_{0}-S|\leq |V_{0}|/2\\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.

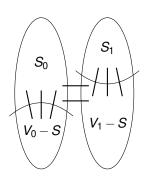


$$\begin{array}{l} \mathcal{S}_{0}|\geq |V_{0}|/2.\\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2\\ \mathcal{S}_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2.\\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}.\\ |S_{0}|\geq |V_{0}|/2\implies |V_{0}-S|\leq |V_{0}|/2\\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



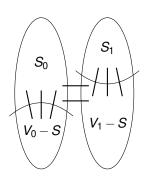
$$\begin{array}{l} S_{0}|\geq |V_{0}|/2.\\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2\\ S_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2.\\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}.\\ |S_{0}|\geq |V_{0}|/2\implies |V_{0}-S|\leq |V_{0}|/2\\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{array}{l} S_{0}|\geq |V_{0}|/2. \\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2 \\ S_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2. \\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}. \\ |S_{0}|\geq |V_{0}|/2 \implies |V_{0}-S|\leq |V_{0}|/2 \\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

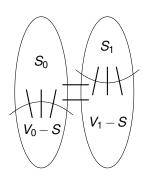
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

>

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{array}{l} S_{0}|\geq |V_{0}|/2.\\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2\\ S_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2.\\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}.\\ |S_{0}|\geq |V_{0}|/2\implies |V_{0}-S|\leq |V_{0}|/2\\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

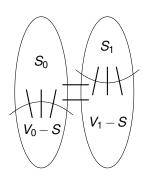
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1|$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{array}{l} S_{0}|\geq |V_{0}|/2.\\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2\\ S_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2.\\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}.\\ |S_{0}|\geq |V_{0}|/2\implies |V_{0}-S|\leq |V_{0}|/2\\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

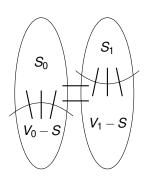
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0|$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{array}{l} S_{0}|\geq |V_{0}|/2.\\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2\\ S_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2.\\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}.\\ |S_{0}|\geq |V_{0}|/2\implies |V_{0}-S|\leq |V_{0}|/2\\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

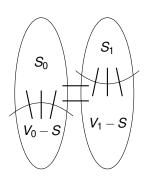
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1|$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{array}{l} S_{0}|\geq |V_{0}|/2. \\ \text{Recall Case 1: } |S_{0}|, |S_{1}|\leq |V|/2 \\ S_{1}|\leq |V_{1}|/2 \text{ since } |S|\leq |V|/2. \\ \implies\geq |S_{1}| \text{ edges cut in } E_{1}. \\ |S_{0}|\geq |V_{0}|/2 \implies |V_{0}-S|\leq |V_{0}|/2 \\ \implies\geq |V_{0}|-|S_{0}| \text{ edges cut in } E_{0}. \end{array}$$

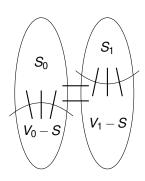
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1| = |V_0|$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| &\leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| &\leq |V|/2. \\ &\implies &\geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| &\leq |V_0|/2 \\ &\implies &\geq |V_0| - |S_0| \text{ edges cut in } E_0. \end{split}$$

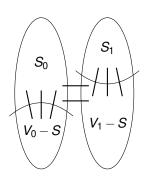
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1| = |V_0|$ $|V_0|$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| \leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| \leq |V|/2. \\ &\implies \geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| \leq |V_0|/2 \\ &\implies \geq |V_0| - |S_0| \text{ edges cut in } E_0. \end{split}$$

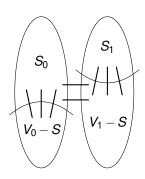
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1| = |V_0| \\ |V_0| = |V|/2 \geq |S|.$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



$$\begin{split} |S_0| &\geq |V_0|/2. \\ \text{Recall Case 1: } |S_0|, |S_1| \leq |V|/2 \\ |S_1| &\leq |V_1|/2 \text{ since } |S| \leq |V|/2. \\ &\implies \geq |S_1| \text{ edges cut in } E_1. \\ |S_0| &\geq |V_0|/2 \implies |V_0 - S| \leq |V_0|/2 \\ &\implies \geq |V_0| - |S_0| \text{ edges cut in } E_0. \end{split}$$

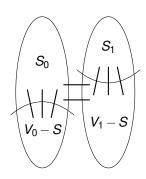
Edges in E_x connect corresponding nodes. $\implies = |S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1| = |V_0| \\ |V_0| = |V|/2 \geq |S|.$

Thm: For any cut (S, V - S) in the hypercube, the number of cut edges is at least the size of the small side, |S|.

Proof: Induction Step. Case 2.



 $|S_0| \ge |V_0|/2.$ Recall Case 1: $|S_0|, |S_1| < |V|/2$ $|S_1| < |V_1|/2$ since |S| < |V|/2. \implies > $|S_1|$ edges cut in E_1 . $|S_0| \ge |V_0|/2 \implies |V_0 - S| \le |V_0|/2$ \implies > $|V_0| - |S_0|$ edges cut in E_0 .

Edges in E_x connect corresponding nodes. \implies = $|S_0| - |S_1|$ edges cut in E_x .

Total edges cut:

 $\geq |S_1| + |V_0| - |S_0| + |S_0| - |S_1| = |V_0|$ $|V_0| = |V|/2 > |S|.$

Also, case 3 where $|S_1| > |V|/2$ is symmetric.

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

Central area of study in computer science!

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

Central area of study in computer science!

Yes/No Computer Programs \equiv Boolean function on $\{0,1\}^n$

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

Central area of study in computer science!

Yes/No Computer Programs \equiv Boolean function on $\{0,1\}^n$

Central object of study.

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

Central area of study in computer science!

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

Central area of study in computer science!

Yes/No Computer Programs \equiv Boolean function on $\{0,1\}^n$

The cuts in the hypercubes are exactly the transitions from 0 sets to 1 set on boolean functions on $\{0,1\}^n$.

Central area of study in computer science!

Yes/No Computer Programs \equiv Boolean function on $\{0,1\}^n$

Central object of study.

Modular Arithmetic.

Applications: cryptography, error correction.

Theorem: If d|x and d|y, then d|(y-x).

Theorem: If d|x and d|y, then d|(y-x). Proof:

Theorem: If d|x and d|y, then d|(y-x).

Proof:

x = ad, y = bd,

Theorem: If d|x and d|y, then d|(y-x).

Proof: x = ad, y = bd, $(x - y) = (ad - bd) = d(a - b) \implies d|(x - y).$

Theorem: If d|x and d|y, then d|(y-x).

Proof: x = ad, y = bd, $(x-y) = (ad-bd) = d(a-b) \implies d|(x-y).$

Theorem: Every number $n \ge 2$ can be represented as a product of primes.

Theorem: If d|x and d|y, then d|(y-x).

Proof: x = ad, y = bd, $(x-y) = (ad-bd) = d(a-b) \implies d|(x-y).$

Theorem: Every number $n \ge 2$ can be represented as a product of primes.

Proof: Either prime, or $n = a \times b$, and use strong induction. (Uniqueness? Later.)

Next Up.

Modular Arithmetic.

If it is 1:00 now.

If it is 1:00 now. What time is it in 2 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours?

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00!

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00. $101 = 12 \times 8 + 5$.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system.

Clock time equivalent up to addition of any integer multiple of 12.

Custom is only to use the representative in $\{12, 1, \dots, 11\}$

If it is 1:00 now. What time is it in 2 hours? 3:00! What time is it in 5 hours? 6:00! What time is it in 15 hours? 16:00! Actually 4:00.

16 is the "same as 4" with respect to a 12 hour clock system. Clock time equivalent up to to addition/subtraction of 12.

What time is it in 100 hours? 101:00! or 5:00.

 $101 = 12 \times 8 + 5.$

5 is the same as 101 for a 12 hour clock system. Clock time equivalent up to addition of any integer multiple of 12.

Custom is only to use the representative in $\{12, 1, ..., 11\}$ (Almost remainder, except for 12 and 0 are equivalent.)

Today is Thursday.

Today is Thursday. What day is it a year from now?

Today is Thursday.

What day is it a year from now? on September 17, 2021?

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4. 5 days from now.

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

25 days from now.

Today is Thursday. What day is it a year from now? on September 17, 2021? Number days. 0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

25 days from now. day 29

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

25 days from now. day 29 or day 1. 29 = (7)4 + 1

two days are equivalent up to addition/subtraction of multiple of 7.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

25 days from now. day 29 or day 1. 29 = (7)4 + 1

two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

25 days from now. day 29 or day 1. 29 = (7)4 + 1

two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday.

25 days from now. day 29 or day 1. 29 = (7)4 + 1

two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

Today is Thursday.
What day is it a year from now? on September 17, 2021?
Number days.
0 for Sunday, 1 for Monday, ..., 6 for Saturday.
Today: day 4.
5 days from now. day 9 or day 2 or Tuesday.
25 days from now. day 29 or day 1. 29 = (7)4 + 1
two days are equivalent up to addition/subtraction of multiple of 7.
11 days from now is day 1 which is Monday!

What day is it a year from now?

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year. So 365 days from now.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year. So 365 days from now. Day 4+365 or day 369.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year. So 365 days from now. Day 4+365 or day 369. Smallest representation:

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year. So 365 days from now. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year. So 365 days from now. Day 4+365 or day 369. Smallest representation:

subtract 7 until smaller than 7. divide and get remainder.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now? Next year is not a leap year. So 365 days from now. Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7. divide and get remainder. 369/7

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now?

Next year is not a leap year. So 365 days from now.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 3.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now?

Next year is not a leap year. So 365 days from now.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 3. 369 = 7(52) + 5

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now?

Next year is not a leap year. So 365 days from now.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 3. 369 = 7(52) + 5

or September 18, 2020 is a Friday.

Today is Thursday.

What day is it a year from now? on September 17, 2021? Number days.

0 for Sunday, 1 for Monday, ..., 6 for Saturday.

Today: day 4.

5 days from now. day 9 or day 2 or Tuesday. 25 days from now. day 29 or day 1. 29 = (7)4 + 1two days are equivalent up to addition/subtraction of multiple of 7. 11 days from now is day 1 which is Monday!

What day is it a year from now?

Next year is not a leap year. So 365 days from now.

Day 4+365 or day 369.

Smallest representation:

subtract 7 until smaller than 7.

divide and get remainder.

369/7 leaves quotient of 52 and remainder 3. 369 = 7(52) + 5

or September 18, 2020 is a Friday.

80 years from now?

80 years from now? 20 leap years.

80 years from now? 20 leap years. 366×20 days

80 years from now? 20 leap years. 366×20 days 60 regular years.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7?

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7?

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4 + 2 \times 20 + 1 \times 60$

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$.

What is remainder of 365 when dividing by 7? 1

Today is day 4.

```
Get Day: 4+2 \times 20+1 \times 60 = 104
Remainder when dividing by 7?
```

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4. Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7$

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4. Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7 + 6$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4. Get Day: $4 + 2 \times 20 + 1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7 + 6$. Or September 18, 2100 is Saturday!

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1 Today is day 4. Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7 + 6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7+6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52\times7+2.$ What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7+6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7+6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7. 60 has remainder 4 when divided by 7. Get Day: $2+2\times 6+1\times 4=18$.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7+6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7. 60 has remainder 4 when divided by 7. Get Day: $2+2\times 6+1\times 4=18$. Or Day 6.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7+6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

Or Day 6. September 18, 2100 is Saturday.

80 years from now? 20 leap years. 366×20 days 60 regular years. 365×60 days Today is day 4. It is day $4 + 366 \times 20 + 365 \times 60$. Equivalent to?

Hmm.

What is remainder of 366 when dividing by 7? $52 \times 7 + 2$. What is remainder of 365 when dividing by 7? 1

Today is day 4.

Get Day: $4+2 \times 20+1 \times 60 = 104$ Remainder when dividing by 7? $104 = 14 \times 7+6$. Or September 18, 2100 is Saturday!

Further Simplify Calculation:

20 has remainder 6 when divided by 7.

60 has remainder 4 when divided by 7.

Get Day: $2 + 2 \times 6 + 1 \times 4 = 18$.

Or Day 6. September 18, 2100 is Saturday.

"Reduce" at any time in calculation!

x is congruent to y modulo m or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by m.

x is congruent to y modulo m or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by m.

... or x and y have the same remainder w.r.t. m.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*. Mod 7 equivalence classes:

Niod / equivalence classes

 $\{\ldots, -7, 0, 7, 14, \ldots\}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*. Mod 7 equivalence classes:

 $\{\dots, -7, 0, 7, 14, \dots\} \ \{\dots, -6, 1, 8, 15, \dots\} \dots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent x and y.

or " $a \equiv c \pmod{m}$ and $b \equiv d \pmod{m}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$ "

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer *k*.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore,

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a + b = c + d + (k + j)m and since k + j is integer.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$?

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m and since k+j is integer. $\implies a+b \equiv c+d \pmod{m}$.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$ "

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m and since k+j is integer. $\implies a+b \equiv c+d \pmod{m}$.

x is congruent to *y* modulo *m* or " $x \equiv y \pmod{m}$ " if and only if (x - y) is divisible by *m*. ...or *x* and *y* have the same remainder w.r.t. *m*. ...or x = y + km for some integer *k*.

Mod 7 equivalence classes:

 $\{\ldots,-7,0,7,14,\ldots\} \ \{\ldots,-6,1,8,15,\ldots\} \ \ldots$

Useful Fact: Addition, subtraction, multiplication can be done with any equivalent *x* and *y*.

or "
$$a \equiv c \pmod{m}$$
 and $b \equiv d \pmod{m}$
 $\implies a+b \equiv c+d \pmod{m}$ and $a \cdot b = c \cdot d \pmod{m}$ "

Proof: If $a \equiv c \pmod{m}$, then a = c + km for some integer k. If $b \equiv d \pmod{m}$, then b = d + jm for some integer j. Therefore, a+b = c+d+(k+j)m and since k+j is integer. $\implies a+b \equiv c+d \pmod{m}$.

Can calculate with representative in $\{0, \ldots, m-1\}$.

x (mod m) or mod(x,m)

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by m in $\{0, \ldots, m-1\}$.

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by m in $\{0, \ldots, m-1\}$.

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by m in $\{0, \ldots, m-1\}$.

 $mod(x,m) = x - \lfloor \frac{x}{m} \rfloor m$

 $x \pmod{m}$ or mod(x,m)- remainder of x divided by $m in \{0, ..., m-1\}$.

mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient.

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient.

 $mod (29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12$

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = 4

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = ¥ = 5

x (mod m) or mod (x,m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system.

x (mod *m*) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29, 12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = \cancel{x} = 5 Work in this system. $a \equiv b \pmod{m}$.

x (mod m) or mod (x,m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system. $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system. 2 = b (mod m)

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system. $a = b \pmod{m}$

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

6 ≡

x (mod m) or mod (x, m)- remainder of x divided by m in $\{0, \ldots, m-1\}$. $mod(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\left|\frac{x}{m}\right|$ is quotient. $mod(29, 12) = 29 - (|\frac{29}{12}|) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers a and b are equivalent modulo m.

Modulus is m

 $6 \equiv 3 + 3$

x (mod m) or mod (x, m)- remainder of x divided by m in $\{0, \ldots, m-1\}$. $mod(x,m) = x - \lfloor \frac{x}{m} \rfloor m$ $\left|\frac{x}{m}\right|$ is quotient. $mod(29, 12) = 29 - (|\frac{29}{12}|) \times 12 = 29 - (2) \times 12 = \cancel{4} = 5$ Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers a and b are equivalent modulo m.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10$

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = X = 5$ Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$.

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system. $a = b \pmod{m}$

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

$$\begin{split} & 6 \equiv 3+3 \equiv 3+10 \pmod{7}. \\ & 6 = \end{split}$$

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$. $6 \equiv 3 + 3$

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system. $2 = b \pmod{m}$

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

- $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$.
- 6 = 3 + 3 = 3 + 10

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = 4 = 5$ Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$.

 $6 = 3 + 3 = 3 + 10 \pmod{7}$.

x (mod m) or mod (x, m) - remainder of x divided by m in {0,...,m-1}. mod (x, m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29, 12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$.

 $6 = 3 + 3 = 3 + 10 \pmod{7}$.

Generally, not 6 $(mod 7) = 13 \pmod{7}$.

x (mod m) or mod (x, m) - remainder of x divided by m in $\{0, ..., m-1\}$. mod $(x, m) = x - \lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod $(29, 12) = 29 - (\lfloor \frac{29}{12} \rfloor) \times 12 = 29 - (2) \times 12 = X = 5$ Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$.

 $6 = 3 + 3 = 3 + 10 \pmod{7}$.

Generally, not 6 $(mod 7) = 13 \pmod{7}$. But probably won't take off points,

x (mod m) or mod (x,m) - remainder of x divided by m in {0,...,m-1}. mod (x,m) = x - $\lfloor \frac{x}{m} \rfloor m$ $\lfloor \frac{x}{m} \rfloor$ is quotient. mod (29,12) = 29 - ($\lfloor \frac{29}{12} \rfloor$) × 12 = 29 - (2) × 12 = X = 5Work in this system.

 $a \equiv b \pmod{m}$.

Says two integers *a* and *b* are equivalent modulo *m*.

Modulus is m

 $6 \equiv 3 + 3 \equiv 3 + 10 \pmod{7}$.

 $6 = 3 + 3 = 3 + 10 \pmod{7}$.

Generally, not 6 $(\mod 7) = 13 \pmod{7}$.

But probably won't take off points, still hard for us to read.

Inverses and Factors.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1;

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$. $2 \cdot 4x = 2 \cdot 5 \pmod{7}$.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

Can solve
$$4x = 5 \pmod{7}$$
.
 $2 \cdot 4x = 2 \cdot 5 \pmod{7}$
 $8x = 10 \pmod{7}$

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

Can solve
$$4x = 5 \pmod{7}$$
.
 $2 \cdot 4x = 2 \cdot 5 \pmod{7}$
 $8x = 10 \pmod{7}$
 $x = 3 \pmod{7}$

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

```
Can solve 4x = 5 \pmod{7}.

2 \cdot 4x = 2 \cdot 5 \pmod{7}

8x = 10 \pmod{7}

x = 3 \pmod{7}

Check!
```

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

Can solve
$$4x = 5 \pmod{7}$$
.
 $2 \cdot 4x = 2 \cdot 5 \pmod{7}$
 $8x = 10 \pmod{7}$
 $x = 3 \pmod{7}$
Check! $4(3) = 12 = 5 \pmod{7}$.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$. $x = 3 \pmod{7}$::: Check! $4(3) = 12 = 5 \pmod{7}$.

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$. $x = 3 \pmod{7}$::: Check! $4(3) = 12 = 5 \pmod{7}$.

For 8 modulo 12: no multiplicative inverse!

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$. $x = 3 \pmod{7}$::: Check! $4(3) = 12 = 5 \pmod{7}$.

For 8 modulo 12: no multiplicative inverse!

"Common factor of 4"

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$. $x = 3 \pmod{7}$::: Check! $4(3) = 12 = 5 \pmod{7}$.

For 8 modulo 12: no multiplicative inverse!

"Common factor of 4" \implies 8*k* - 12*l* is a multiple of four for any *l* and *k* \implies

Division: multiply by multiplicative inverse.

$$2x = 3 \implies \left(\frac{1}{2}\right) \cdot 2x = \left(\frac{1}{2}\right) \cdot 3 \implies x = \frac{3}{2}.$$

Multiplicative inverse of x is y where xy = 1; 1 is multiplicative identity element.

In modular arithmetic, 1 is the multiplicative identity element.

Multiplicative inverse of $x \mod m$ is y with $xy = 1 \pmod{m}$.

For 4 modulo 7 inverse is 2: $2 \cdot 4 \equiv 8 \equiv 1 \pmod{7}$.

Can solve $4x = 5 \pmod{7}$. $x = 3 \pmod{7}$::: Check! $4(3) = 12 = 5 \pmod{7}$.

For 8 modulo 12: no multiplicative inverse!

"Common factor of 4" \implies 8k - 12l is a multiple of four for any l and k \implies 8k \neq 1 (mod 12) for any k.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim:

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$,

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k. gcd(x,m) = 1

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

 \implies Prime factorization of *m* and *x* do not contain common primes.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

 \implies Prime factorization of *m* and *x* do not contain common primes.

 \implies (a-b) factorization contains all primes in *m*'s factorization.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

$$\implies (a-b) \ge m.$$

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

 \implies $(a-b) \ge m$. But $a, b \in \{0, \dots m-1\}$.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

 \implies $(a-b) \ge m$. But $a, b \in \{0, ..., m-1\}$. Contradiction.

Thm:

If greatest common divisor of x and m, gcd(x, m), is 1, then x has a multiplicative inverse modulo m.

Proof \implies : **Claim:** The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Each of m numbers in S correspond to different one of m equivalence classes modulo m.

 \implies One must correspond to 1 modulo *m*. Inverse Exists!

Proof of Claim: If not distinct, then $\exists a, b \in \{0, ..., m-1\}$, $a \neq b$, where $(ax \equiv bx \pmod{m}) \implies (a-b)x \equiv 0 \pmod{m}$ Or (a-b)x = km for some integer k.

gcd(x,m) = 1

⇒ Prime factorization of *m* and *x* do not contain common primes. ⇒ (a-b) factorization contains all primes in *m*'s factorization. So (a-b) has to be multiple of *m*.

 \implies $(a-b) \ge m$. But $a, b \in \{0, ..., m-1\}$. Contradiction.

Proof review. Consequence.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof review. Consequence.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains

 $y \equiv 1 \mod m$ if all distinct modulo m.

Proof review. Consequence.

. . .

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4...

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... S =

...

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$

...

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6)

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2. Can't be 1.

...

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, \dots, (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo m. ... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2. Can't be 1. No inverse. For x = 5 and m = 6. $S = \{0, 4, 2, 0, 4, 2\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For
$$x = 5$$
 and $m = 6$.
 $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For
$$x = 5$$
 and $m = 6$.
 $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For
$$x = 5$$
 and $m = 6$.
 $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct,

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1!

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$ Not distinct. Common factor 2. Can't be 1. No inverse

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6).

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6.

 $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6).

(Hmm. What normal number is it own multiplicative inverse?)

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6.

 $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6).

(Hmm. What normal number is it own multiplicative inverse?) 1

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6.

 $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$

All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x?

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

$$5x = 3 \pmod{6}$$
 What is x? Multiply both sides by 5.
x = 15

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

... For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd.

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd. $4x = 2 \pmod{6}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd. $4x = 2 \pmod{6}$ Two solutions!

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd. $4x = 2 \pmod{6}$ Two solutions! $x = 2,5 \pmod{6}$

Thm: If gcd(x, m) = 1, then x has a multiplicative inverse modulo m.

Proof Sketch: The set $S = \{0x, 1x, ..., (m-1)x\}$ contains $y \equiv 1 \mod m$ if all distinct modulo *m*.

For x = 4 and m = 6. All products of 4... $S = \{0(4), 1(4), 2(4), 3(4), 4(4), 5(4)\} = \{0, 4, 8, 12, 16, 20\}$ reducing (mod 6) $S = \{0, 4, 2, 0, 4, 2\}$

Not distinct. Common factor 2. Can't be 1. No inverse.

For x = 5 and m = 6. $S = \{0(5), 1(5), 2(5), 3(5), 4(5), 5(5)\} = \{0, 5, 4, 3, 2, 1\}$ All distinct, contains 1! 5 is multiplicative inverse of 5 (mod 6). (Hmm. What normal number is it own multiplicative inverse?) 1 -1.

 $5x = 3 \pmod{6}$ What is x? Multiply both sides by 5. x = $15 = 3 \pmod{6}$

 $4x = 3 \pmod{6}$ No solutions. Can't get an odd. $4x = 2 \pmod{6}$ Two solutions! $x = 2,5 \pmod{6}$

Very different for elements with inverses.

Proof Review 2: Bijections.

If gcd(x,m) = 1.

Proof Review 2: Bijections.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection.

Proof Review 2: Bijections.

If gcd(x,m) = 1.

Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4},$

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4},$

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$ Oh yeah.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}$. Oh yeah. f(0) = 0.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$ Oh yeah. f(0) = 0.

Bijection

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}$. Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}$. Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}$. Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

All the images are distinct. \implies unique pre-image for any image.

x = 2, m = 4.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$ Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

$$x = 2, m = 4.$$

 $f(1) = 2, f(2) = 0, f(3) = 2$

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$ Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

$$x = 2, m = 4.$$

 $f(1) = 2, f(2) = 0, f(3) = 2$
Oh yeah.

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$ Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

$$x = 2, m = 4.$$

 $f(1) = 2, f(2) = 0, f(3) = 2$
Oh yeah. $f(0) = 0.$

If gcd(x,m) = 1. Then the function $f(a) = xa \mod m$ is a bijection. One to one: there is a unique pre-image. Onto: the sizes of the domain and co-domain are the same. x = 3, m = 4. $f(1) = 3(1) = 3 \pmod{4}, f(2) = 6 = 2 \pmod{4}, f(3) = 1 \pmod{3}.$ Oh yeah. f(0) = 0.

Bijection \equiv unique pre-image and same size.

All the images are distinct. \implies unique pre-image for any image.

$$x = 2, m = 4.$$

 $f(1) = 2, f(2) = 0, f(3) = 2$
Oh yeah. $f(0) = 0.$

Not a bijection.

How to find the inverse?

How to find the inverse?

How to find if x has an inverse modulo m?

How to find the inverse?

How to find if x has an inverse modulo m?

Find gcd (x, m).

How to find the inverse?

How to find if x has an inverse modulo m?

Find gcd (x, m). Greater than 1?

How to find the inverse?

How to find if x has an inverse modulo m?

Find gcd (x, m).

Greater than 1? No multiplicative inverse.

How to find the inverse?

How to find if x has an inverse modulo m?

```
Find gcd (x, m).
Greater than 1? No multiplicative inverse.
Equal to 1?
```

How to find the inverse?

How to find if x has an inverse modulo m?

```
Find gcd (x, m).
```

Greater than 1? No multiplicative inverse. Equal to 1? Multiplicative inverse.

How to find the inverse?

How to find if x has an inverse modulo m?

```
Find gcd (x, m).
Greater than 1? No multiplicative inverse.
Equal to 1? Multiplicative inverse.
```

Algorithm:

How to find the inverse?

How to find if x has an inverse modulo m?

Find gcd (x, m). Greater than 1? No multiplicative inverse.

Equal to 1? Mutliplicative inverse.

Algorithm: Try all numbers up to x to see if it divides both x and m.

How to find the inverse?

How to find if x has an inverse modulo m?

Find gcd (x, m).

Greater than 1? No multiplicative inverse.

Equal to 1? Mutliplicative inverse.

Algorithm: Try all numbers up to x to see if it divides both x and m. Very slow.

How to find the inverse?

How to find if x has an inverse modulo m?

Find gcd (x, m).

Greater than 1? No multiplicative inverse.

Equal to 1? Mutliplicative inverse.

Algorithm: Try all numbers up to x to see if it divides both x and m. Very slow.

Next up.

Next up.

Next up. Euclid's Algorithm.

Next up.

Euclid's Algorithm. Runtime.

Next up.

Euclid's Algorithm. Runtime. Euclid's Extended Algorithm.

Does 2 have an inverse mod 8?

Does 2 have an inverse mod 8? No.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9?

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9?

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$.

Does 6 have an inverse mod 9? No.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$.

Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)!

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)!

x has an inverse modulo m if and only if

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)!x has an inverse modulo m if and only if

gcd(x,m) > 1?

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)! x has an inverse modulo m if and only if gcd(x,m) > 1? No.

gcd(x, m) = 1?

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)! x has an inverse modulo m if and only if gcd(x,m) > 1? No.

gcd(x,m) = 1? Yes.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0 + 8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0 + 9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)! x has an inverse modulo m if and only if gcd(x,m) > 1? No.

gcd(x,m) = 1? Yes.

Now what?:

Compute gcd!

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)!*x* has an inverse modulo *m* if and only if

gcd(x,m) > 1? No. gcd(x,m) = 1? Yes.

Now what?:

Compute gcd!

Compute Inverse modulo *m*.

Does 2 have an inverse mod 8? No. Any multiple of 2 is 2 away from 0+8k for any $k \in \mathbb{N}$. Does 2 have an inverse mod 9? Yes. 5 $2(5) = 10 = 1 \mod 9$. Does 6 have an inverse mod 9? No. Any multiple of 6 is 3 away from 0+9k for any $k \in \mathbb{N}$. 3 = gcd(6,9)!*x* has an inverse modulo *m* if and only if

gcd(x,m) > 1? No. gcd(x,m) = 1? Yes.

Now what?:

Compute gcd!

Compute Inverse modulo *m*.

Notation: *d*|*x* means "*d* divides *x*" or

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Fact: If d|x and d|y then d|(x+y) and d|(x-y).

Notation: d|x means "d divides x" or x = kd for some integer k. Fact: If d|x and d|y then d|(x+y) and d|(x-y). Is it a fact?

Notation: d|x means "d divides x" or x = kd for some integer k. Fact: If d|x and d|y then d|(x + y) and d|(x - y). Is it a fact? Yes?

Notation: d|x means "d divides x" or x = kd for some integer k. Fact: If d|x and d|y then d|(x + y) and d|(x - y). Is it a fact? Yes? No?

Notation: d|x means "d divides x" or x = kd for some integer k. Fact: If d|x and d|y then d|(x+y) and d|(x-y). Is it a fact? Yes? No? Proof: d|x and d|y or

Notation: d|x means "d divides x" or x = kd for some integer k. Fact: If d|x and d|y then d|(x + y) and d|(x - y). Is it a fact? Yes? No? Proof: d|x and d|y or $x = \ell d$ and y = kd

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Fact: If d|x and d|y then d|(x+y) and d|(x-y).

Is it a fact? Yes? No?

Proof: d|x and d|y or $x = \ell d$ and y = kd $\implies x - y = kd - \ell d$

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Fact: If d|x and d|y then d|(x+y) and d|(x-y).

Is it a fact? Yes? No?

Proof: d|x and d|y or $x = \ell d$ and y = kd $\implies x - y = kd - \ell d = (k - \ell)d$

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Fact: If d|x and d|y then d|(x+y) and d|(x-y).

Is it a fact? Yes? No?

Proof:
$$d|x$$
 and $d|y$ or
 $x = \ell d$ and $y = kd$
 $\implies x - y = kd - \ell d = (k - \ell)d \implies d|(x - y)$

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Fact: If d|x and d|y then d|(x+y) and d|(x-y).

Is it a fact? Yes? No?

Proof:
$$d|x$$
 and $d|y$ or
 $x = \ell d$ and $y = kd$
 $\implies x - y = kd - \ell d = (k - \ell)d \implies d|(x - y)$

Notation: *d*|*x* means "*d* divides *x*" or

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

 $mod(x,y) = x - \lfloor x/y \rfloor \cdot y$

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

Therefore $d \mod (x, y)$.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And $d \mid y$ since it is in condition.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar. Try this at home.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar. Try this at home.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar. Try this at home.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

⊡ish.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar. Try this at home.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)). **Proof:** *x* and *y* have **same** set of common divisors as *x* and mod(x, y) by Lemma 1 and 2. □ish.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar. Try this at home.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)). **Proof:** *x* and *y* have **same** set of common divisors as *x* and mod (x, y) by Lemma 1 and 2. Same common divisors \implies largest is the same. ⊡ish.

Notation: d|x means "*d* divides *x*" or x = kd for some integer *k*.

Lemma 1: If d|x and d|y then d|y and $d| \mod (x, y)$.

Proof:

$$\begin{array}{lll} \operatorname{mod} (x,y) &=& x - \lfloor x/y \rfloor \cdot y \\ &=& x - \lfloor s \rfloor \cdot y & \text{for integer } s \\ &=& kd - s\ell d & \text{for integers } k, \ell \text{ where } x = kd \text{ and } y = \ell d \\ &=& (k - s\ell)d \end{array}$$

Therefore $d \mod (x, y)$. And d | y since it is in condition.

Lemma 2: If d|y and $d| \mod (x, y)$ then d|y and d|x. **Proof...:** Similar. Try this at home.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)). **Proof:** *x* and *y* have **same** set of common divisors as *x* and mod (x, y) by Lemma 1 and 2. Same common divisors \implies largest is the same. ⊡ish.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)?

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)?

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*"

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest."

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction.Base Case: y = 0, "x divides y and x" \implies "x is common divisor and clearly largest."Induction Step: $mod(x, y) < y \le x$ when $x \ge y$

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest." **Induction Step:** mod $(x, y) < y \le x$ when $x \ge y$ call in line (***) meets conditions plus arguments "smaller"

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest." **Induction Step:** mod $(x, y) < y \le x$ when $x \ge y$ call in line (***) meets conditions plus arguments "smaller"

and by strong induction hypothesis

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest." **Induction Step:** mod $(x, y) < y \le x$ when $x \ge y$ call in line (***) meets conditions plus arguments "smaller"

and by strong induction hypothesis computes gcd(y, mod(x, y))

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest." **Induction Step:** mod $(x, y) < y \le x$ when $x \ge y$ call in line (***) meets conditions plus arguments "smaller" and by strong induction hypothesis computes gcd(*y*, mod (x, y)) which is gcd(*x*, *y*) by GCD Mod Corollary.

GCD Mod Corollary: gcd(x, y) = gcd(y, mod(x, y)).

Hey, what's gcd(7,0)? 7 since 7 divides 7 and 7 divides 0 What's gcd(x,0)? x

```
(define (euclid x y)
  (if (= y 0)
        x
        (euclid y (mod x y)))) ***
```

Theorem: (euclid x y) = gcd(x, y) if $x \ge y$.

Proof: Use Strong Induction. **Base Case:** y = 0, "*x* divides *y* and *x*" \implies "*x* is common divisor and clearly largest." **Induction Step:** mod $(x, y) < y \le x$ when $x \ge y$ call in line (***) meets conditions plus arguments "smaller" and by strong induction hypothesis computes gcd(*y*, mod (x, y)) which is gcd(*x*, *y*) by GCD Mod Corollary.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For $a \equiv b \pmod{N}$, and $c \equiv d \pmod{N}$, $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For $a \equiv b \pmod{N}$, and $c \equiv d \pmod{N}$, $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division?

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division? Multiply by multiplicative inverse.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For $a \equiv b \pmod{N}$, and $c \equiv d \pmod{N}$, $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For $a \equiv b \pmod{N}$, and $c \equiv d \pmod{N}$, $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a+b=c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why?

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a+b=c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, \ldots, N-1\}$.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a + b = c + d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, \dots, N-1\}$. $ax - ay = 0 \pmod{N} \implies a(x - y)$ is a multiple of N.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, \dots, N-1\}$. $ax - ay = 0 \pmod{N} \implies a(x - y)$ is a multiple of N. If gcd(a, N) = 1,

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, ..., N-1\}$. $ax - ay = 0 \pmod{N} \implies a(x - y)$ is a multiple of *N*. If gcd(a, N) = 1, then (x - y) must contain all primes in prime factorization of *N*,

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a + b = c + d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, ..., N-1\}$. $ax - ay = 0 \pmod{N} \implies a(x - y)$ is a multiple of *N*. If gcd(a, N) = 1, then (x - y) must contain all primes in prime factorization of *N*, and is therefore be bigger than *N*.

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a + b = c + d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, \dots, N-1\}$. $ax - ay = 0 \pmod{N} \implies a(x - y)$ is a multiple of *N*. If gcd(a, N) = 1, then (x - y) must contain all primes in prime factorization of *N*, and is therefore be bigger than *N*.

Only if: For a = xd and N = yd,

Modular Arithmetic: $x \equiv y \pmod{N}$ if x = y + kN for some integer k.

For
$$a \equiv b \pmod{N}$$
, and $c \equiv d \pmod{N}$,
 $ac = bd \pmod{N}$ and $a+b = c+d \pmod{N}$.

Division? Multiply by multiplicative inverse. $a \pmod{N}$ has multiplicative inverse, $a^{-1} \pmod{N}$. If and only if gcd(a, N) = 1.

Why? If: $f(x) = ax \pmod{N}$ is a bijection on $\{1, ..., N-1\}$. $ax - ay = 0 \pmod{N} \implies a(x - y)$ is a multiple of *N*. If gcd(a, N) = 1, then (x - y) must contain all primes in prime factorization of *N*, and is therefore be bigger than *N*. Only if: For a = xd and N = yd,

any ma + kN = d(mx - ky) or is a multiple of d, and is not 1.